

# NORX

A Parallel and Scalable Authenticated Encryption Scheme

Jean-Philippe Aumasson<sup>1</sup> (@veorq)

**Philipp Jovanovic**<sup>2</sup> (@daeinar)

Samuel Neves<sup>3</sup> (@sevenps)

<sup>1</sup>Kudelski Security, Switzerland

<sup>2</sup>University of Passau, Germany

<sup>3</sup>University of Coimbra, Portugal

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# Outline

1. Motivation
2. Specification
3. Performance
4. Security Analysis
5. Conclusion

*“Nearly all of the symmetric encryption modes you learned about in school, textbooks, and Wikipedia are (potentially) insecure.”*

—Matthew Green

# Authenticated Encryption

# Authenticated Encryption

## Types

- ▶ **AE**: ensure *confidentiality*, *integrity*, and *authenticity* of a message
- ▶ **AEAD**: AE + ensure *integrity* and *authenticity* of associated data (e.g. routing information in IP packets)

## Generic Composition

- ▶ Symmetric encryption algorithm (confidentiality)
- ▶ Message Authentication Code (MAC) (authenticity, integrity)

## Applications

- ▶ Standard technology to protect in-transit data
- ▶ IPSec, SSH, TLS, ...

# Authenticated Encryption

## Problems with Existing AE(AD) Schemes

- ▶ Interaction flaws between enc. and auth. in generic composition
- ▶ Weak primitives (e.g. RC4)
- ▶ Broken modes (e.g. EAXprime)
- ▶ Misuse resistant solutions barely used
- ▶ No reliable standards
- ▶ More examples: <http://competitions.cr.yp.to/disasters.html>

⇒ Lots of room for improvements

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# CAESAR



- ▶ **Competition for Authenticated Encryption: Security, Applicability, and Robustness.**
- ▶ **Goals:** Identify a portfolio of *authenticated ciphers* that
  - offer advantages over AES-GCM (the current de-facto standard) and
  - are suitable for widespread adoption.
- ▶ **Overview:**
  - March 15 2014 – End of 2017
  - 1st round: 57 submissions
  - <http://competitions.cr.yp.to/caesar.html>
- ▶ **Further Information:**
  - AEZoo: <https://aezoo.compute.dtu.dk>
  - Speed comparison: <http://www1.spms.ntu.edu.sg/~syllab/speed>



# NO(T A)RX

# Overview of NORX

## Main Design Goals

- ▶ High security
- ▶ Efficiency
- ▶ Simplicity
- ▶ Scalability
- ▶ Online
- ▶ Side-channel robustness  
(e.g. constant-time operations)
- ▶ High key agility

# Overview of NORX

## Parameters

- ▶ *Word size:*  $W \in \{32, 64\}$  bits
- ▶ *Number of rounds:*  $1 \leq R \leq 63$
- ▶ *Parallelism degree:*  $0 \leq D \leq 255$
- ▶ *Tag size:*  $|A| \leq 10W$

## Instances

Rank	NORX $W$ - $R$ - $D$	Nonce size ( $2W$ )	Key size ( $4W$ )	Tag size ( $4W$ )	Classification
1	NORX64-4-1	128	256	256	Standard
2	NORX32-4-1	64	128	128	Standard
3	NORX64-6-1	128	256	256	High security
4	NORX32-6-1	64	128	128	High security
5	NORX64-4-4	128	256	256	High throughput

# Overview of NORX

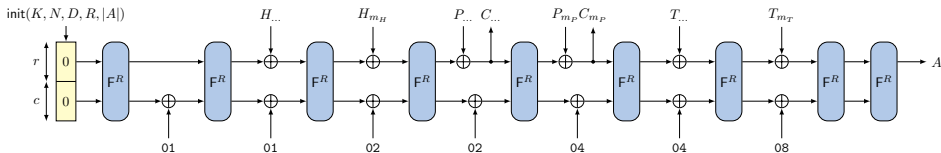
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# NORX Mode

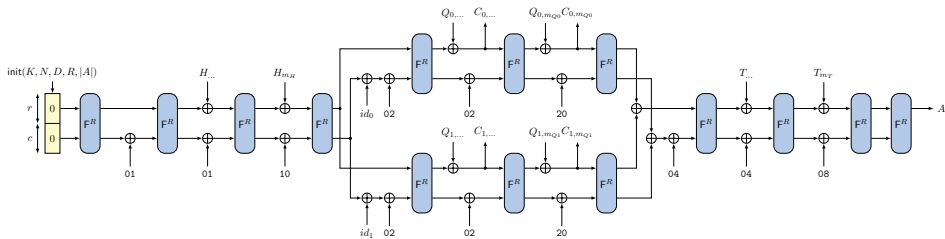


## NORX in Sequential Mode ( $D = 1$ )

### Features

- ▶ (Parallel) monkeyDuplex construction (derived from Keccak/SHA-3)
- ▶ Processes header, payload and trailer data in one-pass
- ▶ Data expansion via multi-rate padding:  $10^*1$
- ▶ Extensible (e.g. sessions, secret message numbers)
- ▶ Parallelisable

# NORX Mode



## NORX in Parallel Mode ( $D = 2$ )

### Features

- ▶ (Parallel) monkeyDuplex construction (derived from Keccak/SHA-3)
- ▶ Processes header, payload and trailer data in one-pass
- ▶ Data expansion via multi-rate padding:  $10^*1$
- ▶ Extensible (e.g. sessions, secret message numbers)
- ▶ Parallelisable

# The State

- ▶ NORX operates on a state of  $16$   $W$ -bit sized words

$W$	Size	Rate	Capacity
32	512	320	192
64	1024	640	384

- ▶ Arrangement of **rate** (data processing) and **capacity** (security) words:

$s_0$	$s_1$	$s_2$	$s_3$
$s_4$	$s_5$	$s_6$	$s_7$
$s_8$	$s_9$	$s_{10}$	$s_{11}$
$s_{12}$	$s_{13}$	$s_{14}$	$s_{15}$

# Initialisation

- ▶ Load **nonce**, **key** and **constants** into state  $S$ :

$u_0$	$n_0$	$n_1$	$u_1$
$k_0$	$k_1$	$k_2$	$k_3$
$u_2$	$u_3$	$u_4$	$u_5$
$u_6$	$u_7$	$u_8$	$u_9$

- ▶ Parameter integration:

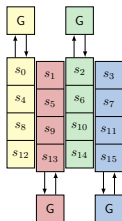
$$s_{14} \leftarrow s_{14} \oplus (R \lll 26) \oplus (D \lll 18) \oplus (W \lll 10) \oplus |A|$$

- ▶ Apply round permutation  $F^R$  to  $S$

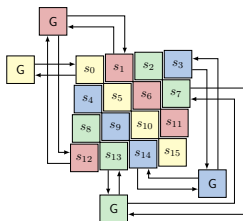


# The Permutation $F^R$

## The Permutation F



## The Permutation G



- 1:  $a \leftarrow H(a, b)$
- 2:  $d \leftarrow (a \oplus d) \ggg r_0$
- 3:  $c \leftarrow H(c, d)$
- 4:  $b \leftarrow (b \oplus c) \ggg r_1$
- 5:  $a \leftarrow H(a, b)$
- 6:  $d \leftarrow (a \oplus d) \ggg r_2$
- 7:  $c \leftarrow H(c, d)$
- 8:  $b \leftarrow (b \oplus c) \ggg r_3$

## The Non-linear Operation H

$$H : \mathbb{F}_2^{2n} \rightarrow \mathbb{F}_2^n, (x, y) \mapsto (x \oplus y) \oplus ((x \wedge y) \ll 1)$$

## Rotation Offsets $(r_0, r_1, r_2, r_3)$

32-bit: (8, 11, 16, 31)

64-bit: (8, 19, 40, 63)

# The Permutation $F^R$

## Features

- ▶ F and G derived from ARX-primitives ChaCha/BLAKE2
- ▶ H is an “approximation” of integer addition:

$$a + b = (a \oplus b) + ((a \wedge b) \ll 1)$$

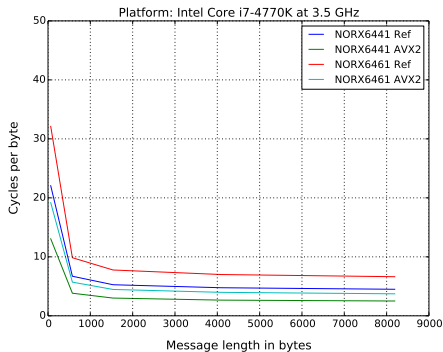
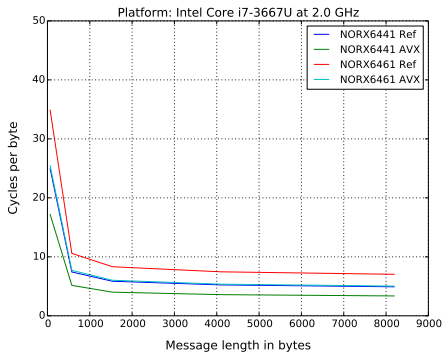
- ▶ LRX permutation
- ▶ No SBoxes or integer additions
- ▶ SIMD friendly
- ▶ HW friendly
- ▶ High diffusion
- ▶ Constant-time

# Requirements for Secure Usage of NORX

1. **Unique nonces**
2. **Abort on tag verification failure**

# Performance

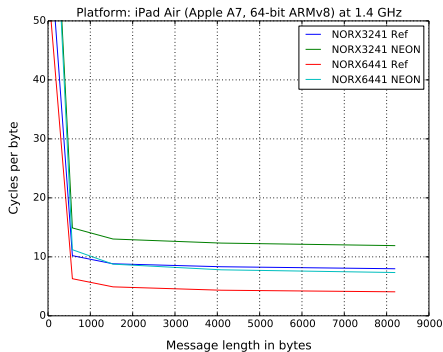
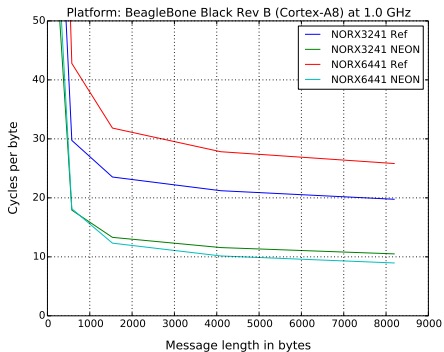
# SW Performance (x86)



Platform	Implementation	cpb	MiBps
Ivy Bridge: i7 3667U @ 2.0 GHz	AVX	3.37	593
Haswell: i7 4770K @ 3.5 GHz	AVX2	2.51	1390

Table: NORX64-4-1 performance

# SW Performance (ARM)

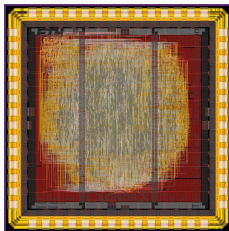


Platform	Implementation	cpb	MiBps
BBB: Cortex-A8 @ 1.0 GHz	NEON	8.96	111
iPad Air: Apple A7 @ 1.4 GHz	Ref	4.07	343

Table: NORX64-4-1 performance



## HW Performance (ASIC)



ASIC implementation and hardware evaluation by ETHZ students (under supervision of Frank K. Gürkaynak):

- ▶ Parameters:  $W \in \{32, 64\}$ ,  $R \in \{2, \dots, 16\}$  and  $D = 1$
- ▶ Technology: 180 nm UMC
- ▶ Frequency: 125 MHz
- ▶ Area requirements: 59 kGE
- ▶ NORX64-4-1 performance: 10 Gbps  $\approx$  1200 MiBps



# Security Analysis

## Sponge Security Bounds

- ▶ Classic:  $\min\{2^{c/2}, 2^{|K|}\}$ 
  - NORX designed towards this bound
  - Expected security levels ( $c - e - 1$ ,  $e = 2W$ ): 127 and 255 bits
- ▶ Improved\*:  $\min\{2^{b/2}, 2^c, 2^{|K|}\}$ 
  - Nonce-based sponges in the ideal perm. model
  - Includes NORX with  $D \neq 1$
  - Effects: rate  $+2W$  bits ( $\approx +16\%$  performance)

\* P. Jovanovic, A. Luykx, and B. Mennink, Beyond  $2^{c/2}$  Security in Sponge-Based Authenticated Encryption Modes, Advances in Cryptology - ASIACRYPT 2014. To appear.

# Cryptanalysis

## NODE – The (NO)RX (D)ifferential Search (E)ngine\*

- ▶ Framework for automatic search of differential trails in  $F^R$
- ▶ Uses constraint / SAT solvers (STP, Boolector, CryptoMiniSat)
- ▶ Some results:

$R$	type	NORX32	NORX64	
1	nonce	$< 2^{-60}$	$< 2^{-53}$	bound
4	perm.	$2^{-584}$	$2^{-836}$	best

- ▶ **Bonus:** Variant of NODE allowed us to break Wheesht and McMambo, two other CAESAR candidates

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Software available on GitHub: <https://github.com/norx/NODE>

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# Conclusion

# Open Problems

- ▶ Cryptanalysis: linear, algebraic, (adv.) differential, (adv.) rotational
- ▶ Side-channel attacks
- ▶ Further implementations: e.g. FPGAs, microcontroller



# Take Aways

## Features of NORX

- ▶ Secure, fast, and scalable
- ▶ Based on well-analysed primitives: ChaCha/BLAKE(2)/Keccak
- ▶ Clean and simple design
- ▶ HW and SW friendly
- ▶ Parallelisable
- ▶ Side-channel robustness considered during design phase
- ▶ Straightforward to implement
- ▶ No AES dependence

## Further Information

<https://norx.io>

Contact:

[jovanovic@fim.uni-passau.de](mailto:jovanovic@fim.uni-passau.de)

@Daeinar

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# Supplement: NORX vs AES-GCM

	NORX	AES-GCM
High performance	yes (on many platforms)	depends (high with AES-NI)
High key agility	yes	no
Timing resistance	yes	no (bit-slicing, AES-NI required)
Misuse resistance	A+N / LCP+X (exposes $P \oplus P'$ )	no (exposes $K$ )
Parallelisation	yes	yes
Extensibility	yes (sessions, secret msg. nr., etc.)	no
Simple implementation	yes	no