## NORX

A Parallel and Scalable Authenticated Encryption Scheme

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Outline

1. Motivation
2. Specification of NORX
3. Analysis of NORX
4. Conclusion

What is Authenticated Encryption?

## Authenticated Encryption?

Tr $\int_{\text {Passau }}^{\text {unversstät }}$
Non-AE



Alice

## Authenticated Encryption?

Non-AE
$C=E_{K}($ Let's meet at 18:00)


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## Authenticated Encryption?

## Non-AE

$C=E_{K}$ (Let's meet at 18:00)
$D_{K}\left(C^{\prime}\right)=$ Let's meet at $20: 00$


## Authenticated Encryption!

Tr $\int_{\text {Passau }}^{\text {unversstät }}$
AE


Bob


Alice

## Authenticated Encryption!

AE
$(C, T)=A E E_{K}($ Let's meet at 18:00)


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## Authenticated Encryption!

AE
$(C, T)=A E E_{K}($ Let's meet at 18:00)

$$
A E D_{K}\left(C^{\prime}, T\right)=\left(P^{\prime}, T^{\prime}\right), T \neq T^{\prime}
$$



## Authenticated Encryption

Types

- AE: ensure confidentiality, integrity, and authenticity of a message.
- AEAD: AE + ensure integrity and authenticity of associated data (e.g. routing information in IP packets).

Generic Composition

- Symmetric encryption algorithm (confidentiality)
- Message Authentication Code (MAC) (integrity)


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## Crypto Disasters

## Problems with Existing AEAD Schemes

- Interaction flaws: enc. $\longleftrightarrow$ auth. (generic composition)
- Weak primitives (e.g. RC4)
- Broken modes (e.g. EAXprime)
- No misuse resistant solutions
- ...
- More examples: http://competitions.cr.yp.to/disasters.html
$\Rightarrow$ Lots of room for improvements


## Crypto Disasters

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$\Rightarrow$ Lots of room for improvements ...
- Competition for Authenticated Encryption: Security, Applicability, and Robustness.
- Goals: Identify a portfolio of authenticated ciphers (one primitive) that
- offer advantages over AES-GCM (the current de-facto standard) and
- are suitable for widespread adoption.
- Overview:
- March 152014 - End of 2017
- 1st round: 57 submissions
- http://competitions.cr.yp.to/caesar.html

Further Information:

- AEZoo: https://aezoo.compute.dtu.dk
- Speed comparison: http://www1.spms.ntu.edu.sg/~syllab/speed


## NORX

## Overview of NORX

## Main Design Goals

- High security
- Efficiency
- Simplicity
- Scalability
- Online
- Single pass
- Timing resistance
- High key agility


## Overview of NORX



## General

- Family of AEAD schemes
- Type: nonce-based stream cipher
> - Mode: (parallel) MonkeyDuplex (introduced with Keccak)
> - Core: LRX permutation (from ChaCha / BLAKE2, ARX-based)
> - Name: "NO(T A)RX"

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## Overview of NORX

Parameters

- Word size: $W \in\{32,64\}$ bits
- Number of rounds: $1 \leq R \leq 63$
- Parallelism degree: $0 \leq D \leq 255$
- Tag size: $|A| \leq 10 W$ (default: $4 W$ bits)
- Input: key K (4W bits), nonce $N(2 W$ bits), and message $M=H\|P\| T$ with $H$ header, $P$ payload, and $T$ trailer.
- Output: encrypted payload C and authentication tag A.


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## Encryption Mode

- Input: key $K$ ( 4 W bits), nonce $N(2 W$ bits), and message $M=H\|P\| T$ with $H$ header, $P$ payload, and $T$ trailer.
- Output: encrypted payload $C$ and authentication tag $A$.


## Overview of NORX

## Proposed Instances

|  | NORXW-R-D | Key size | Tag size | Classification |
| :---: | :---: | :---: | :---: | :---: |
| 1 | NORX64-4-1 | 256 | 256 | standard |
| 2 | NORX32-4-1 | 128 | 128 | standard |
| 3 | NORX64-6-1 | 256 | 256 | high security |
| 4 | NORX32-6-1 | 128 | 128 | high security |
| 5 | NORX64-4-4 | 256 | 256 | high throughput |

arget Patorms

- NORX32: 8- to 32-bit CPUs, low-resource hardware - NORX64: 64-bit CPUs, high performance hardware


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## Target Platforms

- NORX32: 8- to 32-bit CPUs, low-resource hardware
- NORX64: 64-bit CPUs, high performance hardware


## MonkeyDuplex



## The Encryption / Decryption Process



Figure: NORX in Sequential Mode ( $D=1$ )

## The Encryption / Decryption Process



Figure: NORX in Parallel Mode $(D=2)$

## The State

- NORX operates on a state of 16 W -bit sized words

|  | Size | Rate | Capacity |
| :---: | :---: | :---: | :---: |
| NORX32 | 512 | 320 | 192 |
| NORX64 | 1024 | 640 | 384 |

- Arrangement of rate (data processing) and capacity (security) words:

| $s_{0}$ | $s_{1}$ | $s_{2}$ | $s_{3}$ |
| :--- | :--- | :--- | :--- |
| $s_{4}$ | $s_{5}$ | $s_{6}$ | $s_{7}$ |
| $s_{8}$ | $s_{9}$ | $s_{10}$ | $s_{11}$ |
| $s_{12}$ | $s_{13}$ | $s_{14}$ | $s_{15}$ |

## Initialisation

- Load nonce, key and constants into state $S$ :

| $u_{0}$ | $n_{0}$ | $n_{1}$ | $u_{1}$ |
| :---: | :---: | :---: | :---: |
| $k_{0}$ | $k_{1}$ | $k_{2}$ | $k_{3}$ |
| $u_{2}$ | $u_{3}$ | $u_{4}$ | $u_{5}$ |
| $u_{6}$ | $u_{7}$ | $u_{8}$ | $u_{9}$ |

- Parameter integration:

$$
s_{14} \leftarrow s_{14} \oplus(R \ll 26) \oplus(D \ll 18) \oplus(W \ll 10) \oplus|A|
$$

- Apply round permutation $\mathrm{F}^{R}$ to $S$


## The Core Permutation $\mathrm{F}^{R}$

The Permutation F


## The Permutation G

$$
\begin{aligned}
& 1: a \longleftarrow(a \oplus b) \oplus((a \wedge b) \ll 1) \\
& 2: d \longleftarrow(a \oplus d) \ggg r_{0} \\
& 3: c \longleftarrow(c \oplus d) \oplus((c \wedge d) \ll 1) \\
& 4: b \longleftarrow(b \oplus c) \ggg r_{1} \\
& 5: a \longleftarrow(a \oplus b) \oplus((a \wedge b) \ll 1) \\
& 6: d \longleftarrow(a \oplus d) \ggg r_{2} \\
& 7: c \longleftarrow(c \oplus d) \oplus((c \wedge d) \ll 1) \\
& 8: b \longleftarrow(b \oplus c) \ggg r_{3}
\end{aligned}
$$

## Rotation Offsets

- NORX32: $\left(r_{0}, r_{1}, r_{2}, r_{3}\right)=(8,11,16,31)$
$\Rightarrow$ NORX64: $\left(r_{0}, r_{1}, r_{2}, r_{3}\right)=(8,19,40,63)$


## Security Goals

Requirements for secure usage of NORX:

1. Unique nonces
2. Abort on tag verification failure

Expected security levels (in bits):

| Security goal | NORX32 | NORX64 |
| :--- | :---: | :---: |
| Plaintext confidentiality | 128 | 256 |
| Plaintext integrity | 128 | 256 |
| Associated data integrity | 128 | 256 |
| Public message number integrity | 128 | 256 |

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## Sponge Security Level

Classical Bound

$$
\min \left\{2^{c / 2}, 2^{|K|}\right\}
$$

- NORX designed towards this bound
- Usage exponent $e=2 W$, i.e. 64 and 128
- Minimal expected security levels $(c-e-1): 127$ and 255 bits

$$
\min \left\{2^{b / 2}, 2^{c}, 2^{|K|}\right\}
$$

- For nonce-based sponges in the ideal permutation model
- Also includes MORX with $D \neq 1$
- Effects: rate +2 W bits ( $\approx+16 \%$ performance)

[^0]
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## Improved Bound*

$$
\min \left\{2^{b / 2}, 2^{c}, 2^{|K|}\right\}
$$

- For nonce-based sponges in the ideal permutation model
- Also includes NORX with $D \neq 1$
- Effects: rate $+2 W$ bits ( $\approx+16 \%$ performance)
* P. Jovanovic, A. Luykx, and B. Mennink, Beyond $2^{c / 2}$ Security in Sponge-Based Authenticated

Encryption Modes, Cryptology ePrint Archive: Report 2014/373

## Performance of NORX

## SW Performance (x86)




| Platform | Implementation | cpb | MiBps |
| :---: | :---: | :---: | :---: |
| Ivy Bridge: i7 3667U @ 2.0 GHz | AVX | 3.37 | 593 |
| Haswell: i7 4770K @ 3.5 GHz | AVX2 | 2.51 | 1390 |

Table: NORX64-4-1 performance

## SW Performance (ARM)




| Platform | Implementation | cpb | MiBps |
| :---: | :---: | :---: | :---: |
| BBB: Cortex-A8 @ 1.0 GHz | NEON | 8.96 | 111 |
| iPad Air: Apple A7 @ 1.4 GHz | Ref | 4.07 | 343 |

Table: NORX64-4-1 performance

## SW Performance (SUPERCOP)

| Etano | wintermute | sac | hy dra9 | h6, dragon | nopa | misandy | gee | nosdy berery | hlmsps |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 2episi281 | 2es128gemvi | morus540128v1 | Exisi28] | Poratily | morus40128v1 | morasty0128v1 | morusstol28v1 | maxkiativi | morusi280128v1 |
| tiosxinv1 | ece 256 gcmyl | morusi280128v1 | aioxinv1 | mores640128v1 | nora 324 1 vl | morusi280128v1 | morus[290256v] | 7x6+6iv1 | morusi2801256v1 |
| acpis128 | mextediv1 | morusi 280256v1 | pexist28 | -orx 3241 w 1 | haslsiv lov1 | morusi280256v] | morus[280128v] | marx 2241 l ( | hslsiv lov1 |
| teq is256 | mix $6461 \times 1$ | mora $(+1 y 1$ | criv256 | (vorus1280128v1 | $4{ }^{3}$ | muraftaly | haslisivax | morx $6 \times 44 \times 1$ | whicestuv1 |
| Eiamixaq12kv1 | mars $32+1 \times 1$ | 2equal $12 \mathrm{R1}$ | kizanceq $12 \mathrm{kv1}$ | marn 1281256v1 | mrorux 1280128 vi | nare 324 va | maratavi | menn26x 1 |  |
| moraa1280128v1 | marti4/44*1 | minacty | \|hiverv1 |  | manis1280236v1 | narkutivi | bhatsivv1 | manus60128v1 | nara 324.v1 |
| morus128028.6v1 | morus 40128 va | Lesplit 28 | Donsi280128v1 | Estaslov1 |  | \|hatsiviov 1 | [2xon96v1 |  | Weskil281 |
| Elverv. |  | norx324y 1 | Heosy ymeq 128128 vx |  | Pesisi281 | 2es128gemv1 | norx644.4v1 | besl28gcmv1 | hslsivel |
| moras 640128 va | 2sson96v1 | 2ugis256 | feoxysmoq256128vi | Ecgisi28] | Whicoshtv1rar3frlti28 | tiaoxinv1 | nora32flvi | 0ex256cmv1 |  |
| kiasuq $128 \times 1$ | 2seon128v1 |  | kiesocq 128 v 1 | Otsechtv1 mr3fit 1256 | whecshtw/ 1 ar 3 Fri1256 |  | [x+6\|v1 | escon 128 y 1 | serisl2 28 |
| doxaysmeq 12812 Bv 1 | zesi 280ates 1 | 20al 28 Bcanv 1 | morusi280256v1 | mbecohivi nrijirlt 2 b | hislsivv1 | \|whexshlv|mu3fl296 | 2ckisi281 | pib4cipher 128v1 |  |
| deaxysmequ226128v1 | ues1280tpv1 | mors $6444 \times 1$ | feoxyx<l 128128 v 1 | buoxinv1 | -0ntabiv | ecgis 1281 | ackis 128 | pib4cipher256x1 |  |
| 2esi28genavt | 2ec236otrw 1 | 2e256gemv1 | panub50128v1 | Regisi28 | xigic1 28 | Lea236gemv1 | Exeni28v1 | mal 2 Ratrw 1 | -xgi2256 |
| 20236genvi | 2esjambuv1 | 2es128atisw 1 | Pemysect25/12801 | Plecechiv1 ari3ricios6 |  | hictivev1 | pibacipher2\%6vt | zeal28atpv1 | fanotinv1 |
| Hooxyseq128128v1 | 2ec296otry v 1 | Les128otrp v 1 | Eesi2 $\mathrm{g}_{\text {gcmv1 }}$ | bsldivel | ecri2256 |  | pifacipher 128 vl | madsh $512 \mathrm{k} 256 \mathrm{n} 256110256 \mathrm{v1}$ | norx 8444 vi |
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| dooxyscq256128v1 |  | pi32ciphar 128v1 | 6es256gcmv1 | ces128gcmvl | bes256gmv1 | 2egis128 | \|silvers] |  | arxekist |
| monativil |  | 202560trs 1 | pe2566pftv1 | Esslevhiv] | hslsiviviv! | nors $6+44 \times 1$ | iocpole256av] | pi33ecipher256v1 | aiverv1 |
| بca 256 cpffyl | stribob192ri | açambur 1 | amatalv | mo2Sbycmvi | wheceshtv1 arsftr7256 | eck i 2256 | incpolel28av1 | pi32cipher128v1 | whecostev1 arSftr 7256 |
| Ethtel |  | 20256atpu1 | \|sl divkov1 | xara $644 \mathrm{tv1}$ | chat | hsl sivhiv1 | ixepolec128v1 | mayamuvi | yacti 12 kr 2 |
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| max 3241v1 |  | 2won9Gy 1 | Whesutw 1 mi3fr1236 | Was | bar | wreamtov2 | Lesi 28 ceffol | exe2scotrvi | meal 2 Ratrpv 1 |
| Ehoestiv Imr 3 friti28 | lomddha256k128 | pi64cipher256v1 | 2anx 3241v1 | Ebal | bbat | \|sream!0v1 | pi32cipher 128v1 |  | iscpok2562v1 |
| ahoeshtv Imr3frit296 |  | pi64cipher 128 V 1 | Imbitivt | 5624 | bbab | 2son96v1 | pi32cipher256v1 | fradsh2266226n 104140160 v 1 | iscpolel28v1 |
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| mins 6444 y 1 | \|likokeyakv| | \|kisweq128v1 | hislsivel | pos 28 cppthw1 | chos |  | yacs 128 v 2 | medsh 256 k 128 m 95 tuu 128 v 1 | bbal |
| xheocdivimis 3 fist256 | xakcyakv1 | Aritob 192r1 | 2arx $6242 \times 1$ | tha | axs128cpithl | xxami $2 \times 1$ | 2cal 2botpv1 | unds sh2 $256 \mathrm{k} 128 \mathrm{n} 96 \mathrm{tu} 96 \times 1$ | ba3 |
| exan96. 1 | ocankeyakx1 | \|riviuov1 | -xan\%6x1 | fiverv 1 | bir | iercamm $12{ }^{2}$ | Fbi2 |  | -6/2 |
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| misiviliv | fivenkes akv1 | fifedues 288 mofv 1 |  | bas | Exdest1280xbuglen 128v1 | pi6tcipher 2.6 V 1 | cbal | Easkeylkv! | bot |
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| tesfombuy | trivizOv1 | riverkeyakv1 | Estsivhiv1 |  | Exadacs 128 cobtuplen 69 v 1 | silvecr 1 | cbos | tribob192r1 | cbas |
| stram10v2 | 2es128poctv losst | saskeyzkv1 | Erearn10yl | 20udesal2 2 ocblap kes64v1 | bbas | iscpole256av] | cbas | pocankeyakv1 | -bat |
| xxamiov1 |  | pil 6 ciphar096x\| | kramiov2 | kradesal 2 2\%ochty kn9961 | axadacal 92 oxblugicn96v1 | ixpalel28v1 | kbu? | pil fecipher096v1 |  |
| kcpole25fax | raci $28 \mathrm{pactz1g} \mathrm{fl} 128 \mathrm{mal}$ | pil 6 ciphar 128x 1 | 3xan 128v1 | कx ${ }^{4}$ | axadacs1920xblug\|en 64 v 1 | istramı14v2 | kbut | pil 6 ciphar $28 \times 1$ |  |
| kepole 12 avav |  | Lixv11 | cepolel2Riv1 | ,20e128.12 | pexduec1920xtuglen $12 \mathrm{Rv1}$ | iscream14v1 | 2eal 28 n8ctoev1 | riverkegzivi | wexduec $128 \times \times$ tuglen 128.1 |
| xepole 128v1 |  | kelijesv1 | ecpole128.1 | peulues 192 ochtag ken 128 v 1 | thatD | Lees 128 cpforv 1 | 2est28al2clack 1 | ketesal | pect28n8clacv 1 |
| 25005128v1 |  | Ocesinkeyalv 1 | Espole256mv | Pezdeel192ocblafker96v1 | y2es 128v2 | kiaseq 128 l I | Leaddest 12 Sosbtuglen 128v1 | ketejerv1 | 6sal28n12close |
| Escemm12v2 |  | dsoxyserq 128128 v 1 | pi64cipher 256 vl | pezdesi92ocblafker64v1 | norx644.v1 | pi64cipher 128 el | 2eadaes 12 20cblizglen96.1 | Pesal28poetviass | meas6otrw 1 |
| strum 12 v - |  | keticipel | Eramm12v2 | chal0 | peadaes5560cblagien 128v1 | 2scon128v1 | 2eadass 12 2ocbltzelen 6 /v] | 2est28poctll acsi 128 | Sss128cpfbyl |
| Esaream $12 \times 2$ |  |  | erramal 2 v 1 | audes256ocblap kon 128 va \| | exadacs560cblagien 54 vl | \|eses128otrsv | 2est28si22silcv1. | best28pactul [f128mul | esconD6x1 |

Source: http://www1.spms.ntu.edu.sg/~syllab/speed

```
NORX among the fastest CAESAR ciphers
Fastest Sponge-based scheme
Reference implementation has competitive speed, too
```


## SW Performance (SUPERCOP)



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- Fastest Sponge-based scheme
- Reference implementation has competitive speed, too


## HW Performance (ASIC)



ASIC implementation and hardware evaluation by ETHZ students (under supervision of Frank K. Gürkaynak):

- Supported parameters: $W \in\{32,64\}, R \in\{2, \ldots, 16\}$ and $D=1$
- Targeted at high data throughput
- Technology: 180 nm UMC
- Frequency: 125 MHz
- Area requirements: 59 kGE
- NORX64-4-1 performance: $10 \mathrm{Gbps} \approx 1200 \mathrm{MiBps}$


## Analysis of NORX

## Analysis of NORX



The Non-Linear Operation H

$$
H:\{0,1\}^{2 n} \rightarrow\{0,1\}^{n},(a, b) \mapsto(a \oplus b) \oplus((a \wedge b) \ll 1)
$$

Properties

- "Approximation" of integer addition:

$$
a+b=(a \oplus b)+((a \wedge b) \ll 1)
$$

- Carries can only affect the next bit (effects on security?)
- Permutation on $\mathbb{F}_{2}^{n}$ if one input argument is fixed
- Hardware efficiency++
- No SBoxes/integer additions: timing resistance in sw \& hw


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## Analysis of NORX

Diffusion properties of $\mathrm{F}^{R}$ on 1-bit input differences:

|  | NORX32 |  |  |  | ChaCha (32-bit) |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| $R$ | $\min$ | $\max$ | avg | median | $\min$ | $\max$ | avg | median |  |  |  |
| 1 | 83 | 280 | 179.222 | 181 | 73 | 294 | 182.195 | 185 |  |  |  |
| 2 | 194 | 307 | 256.024 | 256 | 199 | 312 | 255.999 | 256 |  |  |  |
| 3 | 198 | 312 | 255.995 | 256 | 204 | 313 | 255.988 | 256 |  |  |  |
| 4 | 201 | 307 | 255.996 | 256 | 200 | 314 | 255.989 | 256 |  |  |  |
|  | NORX64 |  |  |  |  |  |  | ChaCha (64-bit) |  |  |  |
| $R$ | $\min$ | $\max$ | avg | median | $\min$ | max | avg | median |  |  |  |
| 1 | 95 | 429 | 230.136 | 222 | 73 | 506 | 248.843 | 246 |  |  |  |
| 2 | 440 | 589 | 511.982 | 512 | 430 | 591 | 512.013 | 512 |  |  |  |
| 3 | 434 | 589 | 512.008 | 512 | 439 | 589 | 511.971 | 512 |  |  |  |
| 4 | 428 | 589 | 511.986 | 512 | 435 | 585 | 512.008 | 512 |  |  |  |

Full diffusion after $\mathrm{F}^{2}$ (as fast as ChaCha's!)
Diffusion test used in search for non-linear op. / rotation offsets

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| 2 | 194 | 307 | 256.024 | 256 | 199 | 312 | 255.999 | 256 |  |  |  |  |  |  |
| 3 | 198 | 312 | 255.995 | 256 | 204 | 313 | 255.988 | 256 |  |  |  |  |  |  |
| 4 | 201 | 307 | 255.996 | 256 | 200 | 314 | 255.989 | 256 |  |  |  |  |  |  |
|  | NORX64 |  |  |  |  |  |  |  |  |  | ChaCha (64-bit) |  |  |  |
| $R$ | min | max | avg | median | min | max | avg | median |  |  |  |  |  |  |
| 1 | 95 | 429 | 230.136 | 222 | 73 | 506 | 248.843 | 246 |  |  |  |  |  |  |
| 2 | 440 | 589 | 511.982 | 512 | 430 | 591 | 512.013 | 512 |  |  |  |  |  |  |
| 3 | 434 | 589 | 512.008 | 512 | 439 | 589 | 511.971 | 512 |  |  |  |  |  |  |
| 4 | 428 | 589 | 511.986 | 512 | 435 | 585 | 512.008 | 512 |  |  |  |  |  |  |

- Full diffusion after $\mathrm{F}^{2}$ (as fast as ChaCha's!)
- Diffusion test used in search for non-linear op. / rotation offsets

Analysis of NORX
$\stackrel{\text { Tv }}{\text { TrI }} \int_{\text {PASSAU }}^{\text {universität }}$
0. Initialisation

NORX32 Diffusion

2. Diagonal Step


5. Column Step

0. Initialisation

NORX64 Diffusion

2. Diagonal Step


Figure: Diffusion Visualisation of $\mathrm{F}^{R}$.

## Analysis of NORX

## NODE - The (NO)RX (D)ifferential Search (E)ngine*

- Framework for automatic search of differentials in $\mathrm{F}^{R}$
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- Uses constraint / SAT solvers (STP, Boolector, CryptoMiniSat)
- Available on GitHub: https://github.com/norx/NODE
- Best differential trails in $\mathrm{F}^{4}$ (full state):

$$
2^{-584} \text { (32-bit) and } 2^{-836} \text { (64-bit) }
$$

- Differential trail bounds for F (init., diffs in nonce only):
$<2^{-60}$ (32 bit) and $<2^{-53}$ (61 bit)
- Variant of NODE allowed us to break Wheesht and McMambo, two other CAESAR candidates
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## Analysis of NORX

Algebraic Properties of G

|  | \#polynomials by degree |  |  |  |  | \#monomials |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| W | 3 | 4 | 5 | 6 | 7 | 8 | $\min$ | $\max$ | avg | median |
| 32-bit | 2 | 6 | 58 | 2 | 8 | 52 | 12 | 489 | 242 | 49.5 |
| 64-bit | 2 | 6 | 122 | 2 | 8 | 116 | 12 | 489 | 253 | 49.5 |

ANF of F: (direct) construction failed, compute server with 64 GB ran out of memory

High polynomial degree + big number of monomials + large state size: should increase difficulty to mount algebraic attacks

We also examined weak states, fixed points, rotational properties,

## Analysis of NORX

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[^1]
## Analysis of NORX

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## Other Properties of $\mathrm{F}^{R}$

We also examined weak states, fixed points, rotational properties, ...

|  | NORX | AES-GCM |
| :---: | :---: | :---: |
| High performance | yes (on many platforms) | depends (high with AES-NI) |
| High key agility | yes | no |
| Timing resistance | yes | no (bit-slicing, AES-NI required) |
| Misuse resistance | A+N $/$ LCP +X (exposes $P \oplus P^{\prime}$ ) | no (exposes $K$ ) |
| Parallelisation | yes | yes |
| Extensibility | yes (sessions, secret msg. nr., etc.) | no |
| Simple implementation | yes | no |

## Conclusion

- NORX superior to AES-GCM in many important points
- Fast on a broad range of architectures
- Resistance vs timing attacks in hw \& sw (no Int. add. \& no SBoxes)
- Our analysis found no security flaws
- Attacks only on reduced versions / single components
- NORX permutation probably a little bit weaker than ChaCha's
- Additional protection: MonkeyDuplex, restrictive initialisation
- NORX seems to have a good security margin
- However, much more (3rd party (!)) cryptanalysis required
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## Further Information

| $\leftarrow \rightarrow$ Q https://norx.io |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| [NORX] | SPECS \& CODE | FEATURES | CONTACT |  |

$$
[S P E C S \& C O D E]
$$

NORX is not patented and freely available for all applications. Its source code is available under the cCO licence.

A list of 3rd party implementations and results on cryptanalysis can be found here.

## [FEATURES]

https://norx.io

- J-P. Aumasson, P. Jovanovic, and S. Neves, NORX - A First Round Candidate in CAESAR
- J-P. Aumasson, P. Jovanovic, and S. Neves, NORX: Parallel and Scalable AEAD, European Symposium on Research in Computer Security (ESORICS 2014), Wroclaw, Poland
- J-P. Aumasson, P. Jovanovic, and S. Neves, Analysis of NORX, Cryptology ePrint Archive: Report 2014/317
- P. Jovanovic, A. Luykx, and B. Mennink, Beyond $2^{c / 2}$ Security in Sponge-Based Authenticated Encryption Modes, Cryptology ePrint Archive: Report 2014/373


Comic by http://dilbert.com

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https://norx.io


[^0]:    P. Jovanovic, A. Luykx, and B. Mennink, Beyond $2^{c / 2}$ Security in Sponge-Based Authenticated

    Encryption Modes. Cryptology ePrint Archive: Report 2014/373

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