

# NORX

A Parallel and Scalable Authenticated Encryption Scheme

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DTU Compute, Copenhagen  
July 31, 2014

# Outline

1. Motivation
2. Specification of NORX
3. Analysis of NORX
4. Conclusion

# What is Authenticated Encryption?

# Authenticated Encryption?

Non-AE



Bob



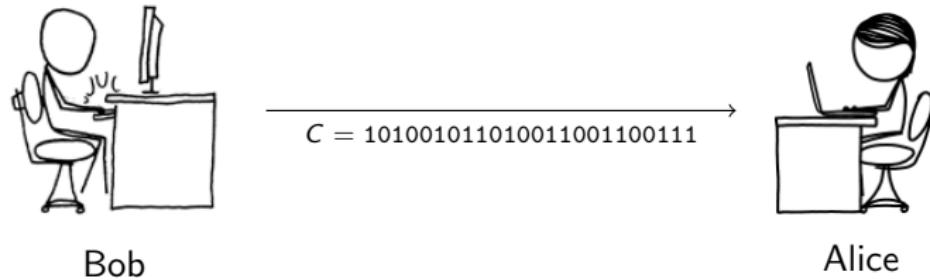
Alice

Images by <https://xkcd.com>

# Authenticated Encryption?

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$C = E_K(\text{Let's meet at 18:00})$

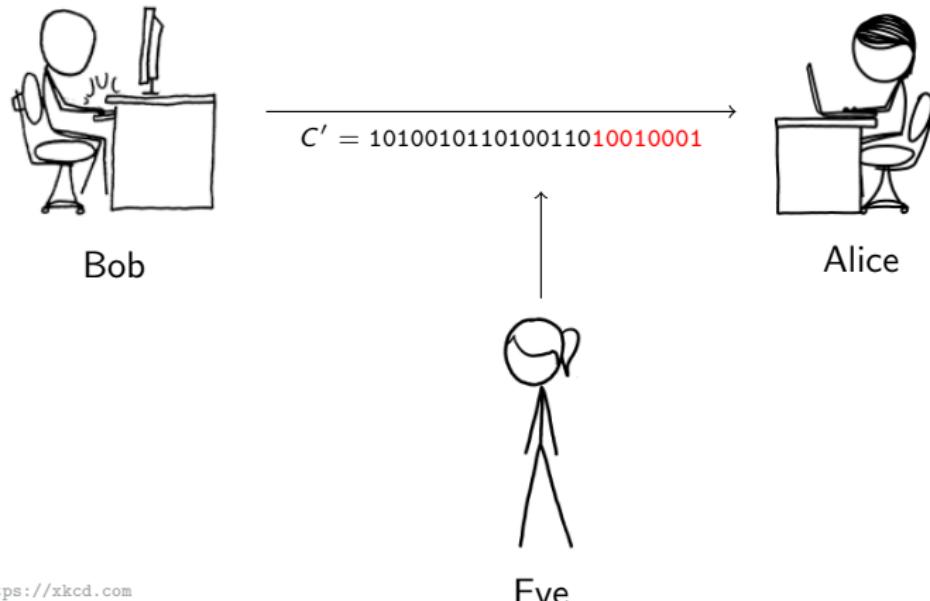


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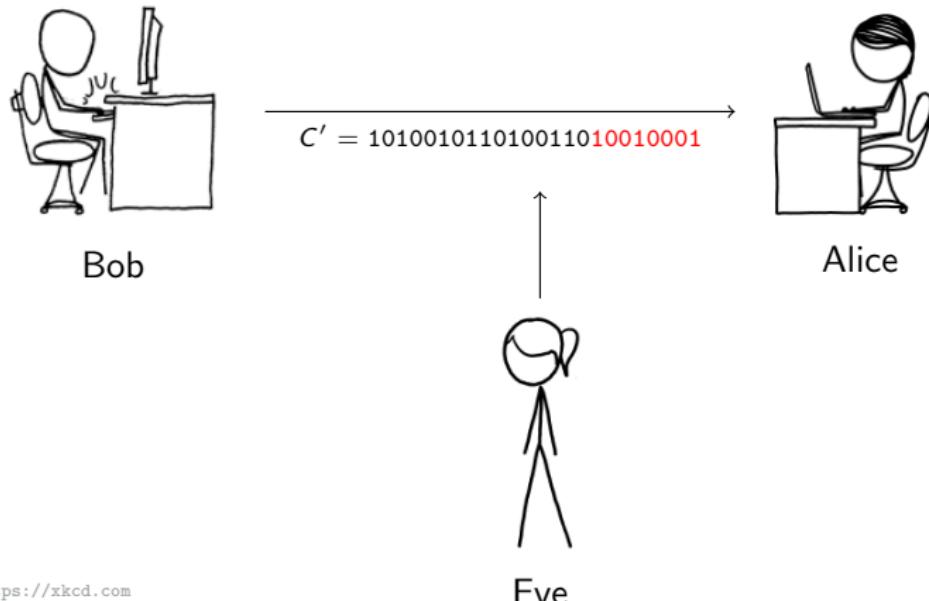
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# Authenticated Encryption?

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$D_K(C') = \text{Let's meet at 20:00}$



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AE



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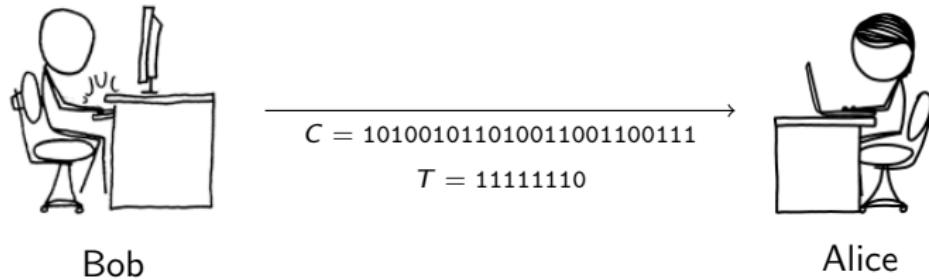
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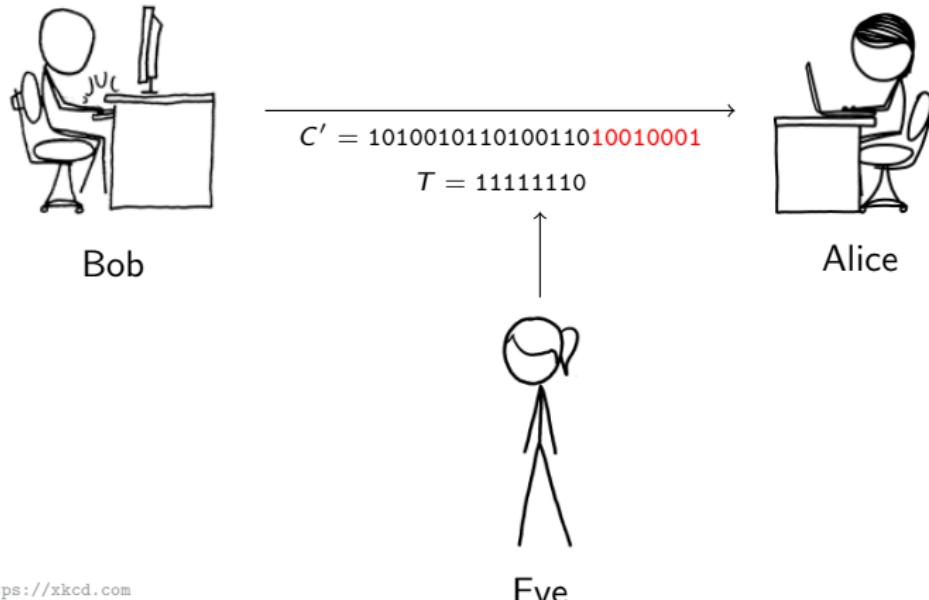


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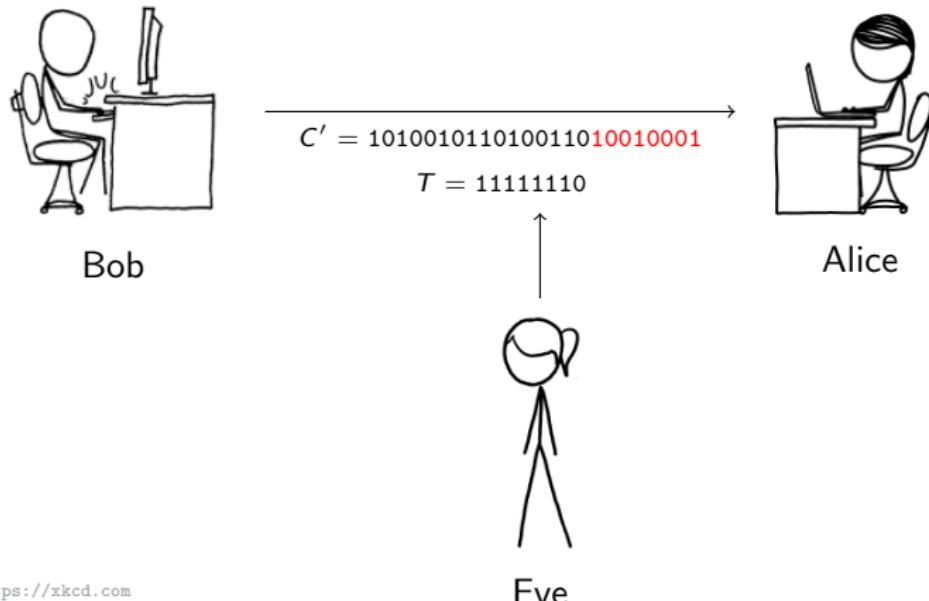
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**AE**

$(C, T) = AEE_K(\text{Let's meet at 18:00})$

$AED_K(C', T) = (P', T')$ ,  $T \neq T'$



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## Types

- ▶ **AE**: ensure *confidentiality*, *integrity*, and *authenticity* of a message.
- ▶ **AEAD**: AE + ensure *integrity* and *authenticity* of associated data (e.g. routing information in IP packets).

## Generic Composition

- ▶ Symmetric encryption algorithm (confidentiality)
- ▶ Message Authentication Code (MAC) (integrity)

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## Problems with Existing AEAD Schemes

- ▶ Interaction flaws: enc.  $\longleftrightarrow$  auth. (generic composition)
- ▶ Weak primitives (e.g. RC4)
- ▶ Broken modes (e.g. EAXprime)
- ▶ No misuse resistant solutions
- ▶ ...
- ▶ More examples: <http://competitions.cr.yp.to/disasters.html>

⇒ Lots of room for improvements . . .

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- ▶ Competition for **A**uthenticated **E**ncryption: **S**ecurity, **A**plicability, and **R**obustness.
- ▶ **Goals:** Identify a portfolio of *authenticated ciphers* (one primitive) that
  - offer advantages over AES-GCM (the current de-facto standard) and
  - are suitable for widespread adoption.
- ▶ **Overview:**
  - March 15 2014 – End of 2017
  - 1st round: 57 submissions
  - <http://competitions.cr.yp.to/caesar.html>
- ▶ **Further Information:**
  - AEZoo: <https://aezoo.compute.dtu.dk>
  - Speed comparison: <http://www1.spms.ntu.edu.sg/~syllab/speed>

# NORX

## Main Design Goals

- ▶ High security
- ▶ Efficiency
- ▶ Simplicity
- ▶ Scalability
- ▶ Online
- ▶ Single pass
- ▶ Timing resistance
- ▶ High key agility

## General

- ▶ Family of AEAD schemes
- ▶ Type: *nonce-based stream cipher*
- ▶ Mode: *(parallel) MonkeyDuplex* (introduced with Keccak)
- ▶ Core: *LRX permutation* (from ChaCha / BLAKE2, ARX-based)
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# Overview of NORX

## Parameters

- ▶ *Word size:*  $W \in \{32, 64\}$  bits
- ▶ *Number of rounds:*  $1 \leq R \leq 63$
- ▶ *Parallelism degree:*  $0 \leq D \leq 255$
- ▶ *Tag size:*  $|A| \leq 10W$  (default:  $4W$  bits)

## Encryption Mode

- ▶ Input: key  $K$  ( $4W$  bits), nonce  $N$  ( $2W$  bits), and message  $M = H \parallel P \parallel T$  with  $H$  header,  $P$  payload, and  $T$  trailer.
- ▶ Output: encrypted payload  $C$  and authentication tag  $A$ .

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- ▶ Output: *encrypted payload C* and *authentication tag A*.

## Proposed Instances

|   | NORXW-R-D  | Key size | Tag size | Classification  |
|---|------------|----------|----------|-----------------|
| 1 | NORX64-4-1 | 256      | 256      | standard        |
| 2 | NORX32-4-1 | 128      | 128      | standard        |
| 3 | NORX64-6-1 | 256      | 256      | high security   |
| 4 | NORX32-6-1 | 128      | 128      | high security   |
| 5 | NORX64-4-4 | 256      | 256      | high throughput |

## Target Platforms

- ▶ NORX32: 8- to 32-bit CPUs, low-resource hardware
- ▶ NORX64: 64-bit CPUs, high performance hardware

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## MonkeyDuplex



# The Encryption / Decryption Process

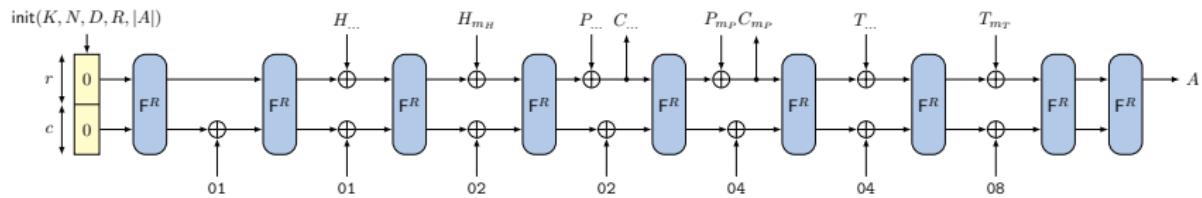


Figure: NORX in Sequential Mode ( $D = 1$ )

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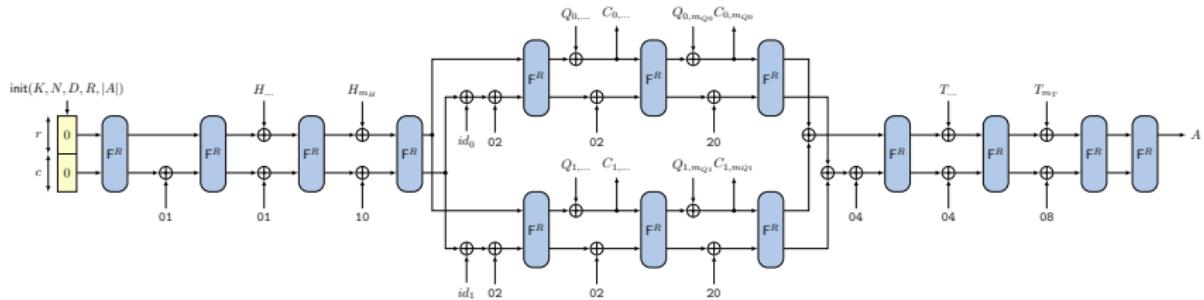


Figure: NORX in Parallel Mode ( $D = 2$ )

# The State

- NORX operates on a state of  $16 W$ -bit sized words

|        | Size | Rate | Capacity |
|--------|------|------|----------|
| NORX32 | 512  | 320  | 192      |
| NORX64 | 1024 | 640  | 384      |

- Arrangement of **rate** (data processing) and **capacity** (security) words:

|          |          |          |          |
|----------|----------|----------|----------|
| $s_0$    | $s_1$    | $s_2$    | $s_3$    |
| $s_4$    | $s_5$    | $s_6$    | $s_7$    |
| $s_8$    | $s_9$    | $s_{10}$ | $s_{11}$ |
| $s_{12}$ | $s_{13}$ | $s_{14}$ | $s_{15}$ |

- ▶ Load **nonce**, **key** and **constants** into state  $S$ :

|       |       |       |       |
|-------|-------|-------|-------|
| $u_0$ | $n_0$ | $n_1$ | $u_1$ |
| $k_0$ | $k_1$ | $k_2$ | $k_3$ |
| $u_2$ | $u_3$ | $u_4$ | $u_5$ |
| $u_6$ | $u_7$ | $u_8$ | $u_9$ |

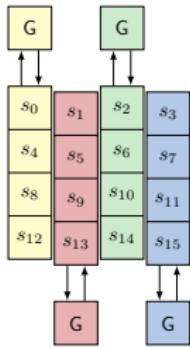
- ▶ Parameter integration:

$$s_{14} \leftarrow s_{14} \oplus (R \ll 26) \oplus (D \ll 18) \oplus (W \ll 10) \oplus |A|$$

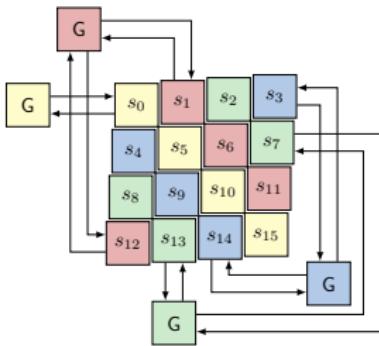
- ▶ Apply round permutation  $F^R$  to  $S$

# The Core Permutation $F^R$

The Permutation F



The Permutation G



- 1 :  $a \leftarrow (a \oplus b) \oplus ((a \wedge b) \ll 1)$
- 2 :  $d \leftarrow (a \oplus d) \ggg r_0$
- 3 :  $c \leftarrow (c \oplus d) \oplus ((c \wedge d) \ll 1)$
- 4 :  $b \leftarrow (b \oplus c) \ggg r_1$
- 5 :  $a \leftarrow (a \oplus b) \oplus ((a \wedge b) \ll 1)$
- 6 :  $d \leftarrow (a \oplus d) \ggg r_2$
- 7 :  $c \leftarrow (c \oplus d) \oplus ((c \wedge d) \ll 1)$
- 8 :  $b \leftarrow (b \oplus c) \ggg r_3$

## Rotation Offsets

- NORX32:  $(r_0, r_1, r_2, r_3) = (8, 11, 16, 31)$
- NORX64:  $(r_0, r_1, r_2, r_3) = (8, 19, 40, 63)$

Requirements for secure usage of NORX:

- 1. Unique nonces**
- 2. Abort on tag verification failure**

Expected security levels (in bits):

| Security goal                   | NORX32 | NORX64 |
|---------------------------------|--------|--------|
| Plaintext confidentiality       | 128    | 256    |
| Plaintext integrity             | 128    | 256    |
| Associated data integrity       | 128    | 256    |
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## Classical Bound

$$\min\{2^{c/2}, 2^{|K|}\}$$

- ▶ NORX designed towards this bound
- ▶ Usage exponent  $e = 2W$ , i.e. 64 and 128
- ▶ Minimal expected security levels ( $c - e - 1$ ): 127 and 255 bits

## Improved Bound\*

$$\min\{2^{b/2}, 2^c, 2^{|K|}\}$$

- ▶ For nonce-based sponges in the ideal permutation model
- ▶ Also includes NORX with  $D \neq 1$
- ▶ Effects: rate  $+2W$  bits ( $\approx +16\%$  performance)

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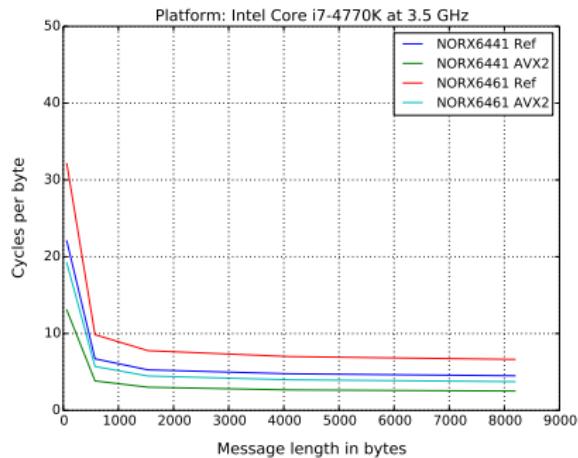
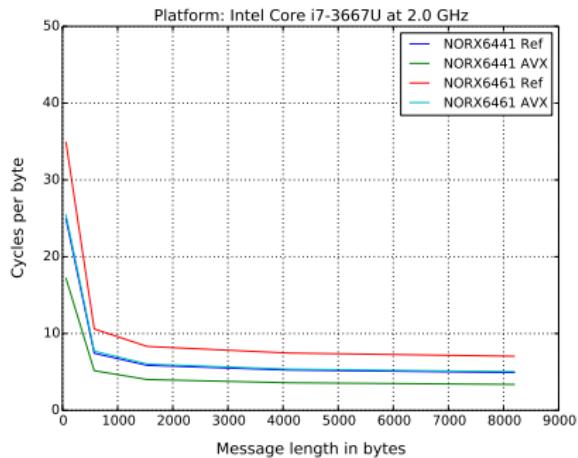
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# Performance of NORX

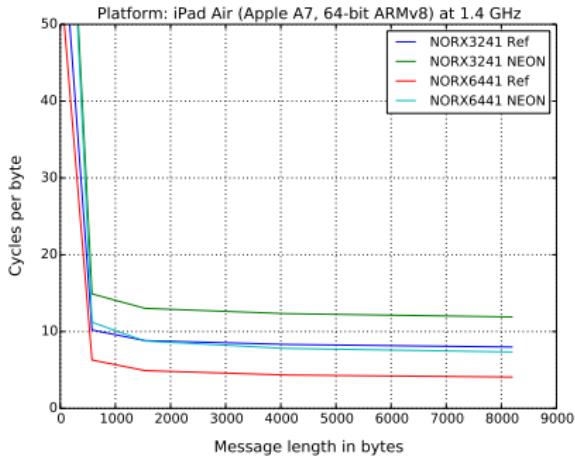
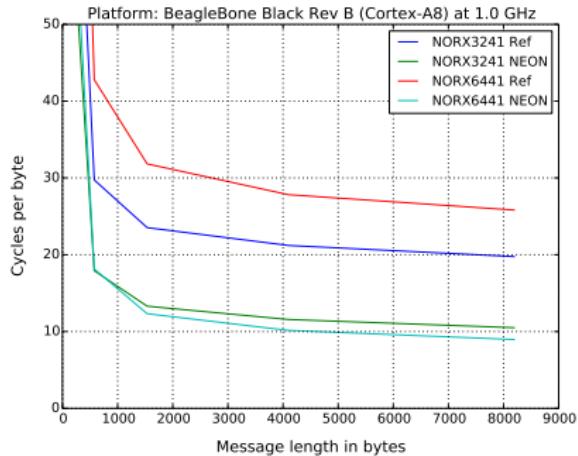
# SW Performance (x86)



| Platform                       | Implementation | cpb  | MiBps |
|--------------------------------|----------------|------|-------|
| Ivy Bridge: i7 3667U @ 2.0 GHz | AVX            | 3.37 | 593   |
| Haswell: i7 4770K @ 3.5 GHz    | AVX2           | 2.51 | 1390  |

Table: NORX64-4-1 performance

# SW Performance (ARM)

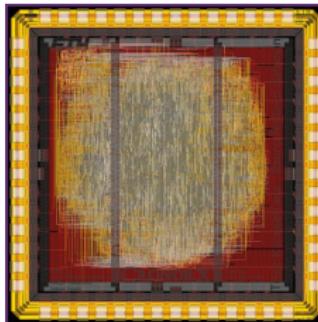


| Platform                     | Implementation | cpb  | MiBps |
|------------------------------|----------------|------|-------|
| BBB: Cortex-A8 @ 1.0 GHz     | NEON           | 8.96 | 111   |
| iPad Air: Apple A7 @ 1.4 GHz | Ref            | 4.07 | 343   |

Table: NORX64-4-1 performance







ASIC implementation and hardware evaluation by ETHZ students (under supervision of Frank K. Gürkaynak):

- ▶ Supported parameters:  $W \in \{32, 64\}$ ,  $R \in \{2, \dots, 16\}$  and  $D = 1$
- ▶ Targeted at high data throughput
- ▶ Technology: 180 nm UMC
- ▶ Frequency: 125 MHz
- ▶ Area requirements: 59 kGE
- ▶ NORX64-4-1 performance: 10 Gbps  $\approx$  1200 MiBps

# Analysis of NORX

## The Non-Linear Operation H

$$H : \{0,1\}^{2n} \rightarrow \{0,1\}^n, (a, b) \mapsto (a \oplus b) \oplus ((a \wedge b) \ll 1)$$

## Properties

- ▶ “Approximation” of integer addition:

$$a + b = (a \oplus b) + ((a \wedge b) \ll 1)$$

- ▶ Carries can only affect the next bit (effects on security?)
- ▶ Permutation on  $\mathbb{F}_2^n$  if one input argument is fixed
- ▶ Hardware efficiency++
- ▶ No SBoxes/integer additions: timing resistance in sw & hw

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Diffusion properties of  $F^R$  on 1-bit input differences:

|     | NORX32 |     |         |        | ChaCha (32-bit) |     |         |        |
|-----|--------|-----|---------|--------|-----------------|-----|---------|--------|
| $R$ | min    | max | avg     | median | min             | max | avg     | median |
| 1   | 83     | 280 | 179.222 | 181    | 73              | 294 | 182.195 | 185    |
| 2   | 194    | 307 | 256.024 | 256    | 199             | 312 | 255.999 | 256    |
| 3   | 198    | 312 | 255.995 | 256    | 204             | 313 | 255.988 | 256    |
| 4   | 201    | 307 | 255.996 | 256    | 200             | 314 | 255.989 | 256    |
|     | NORX64 |     |         |        | ChaCha (64-bit) |     |         |        |
| $R$ | min    | max | avg     | median | min             | max | avg     | median |
| 1   | 95     | 429 | 230.136 | 222    | 73              | 506 | 248.843 | 246    |
| 2   | 440    | 589 | 511.982 | 512    | 430             | 591 | 512.013 | 512    |
| 3   | 434    | 589 | 512.008 | 512    | 439             | 589 | 511.971 | 512    |
| 4   | 428    | 589 | 511.986 | 512    | 435             | 585 | 512.008 | 512    |

- ▶ Full diffusion after  $F^2$  (as fast as ChaCha's!)
- ▶ Diffusion test used in search for non-linear op. / rotation offsets

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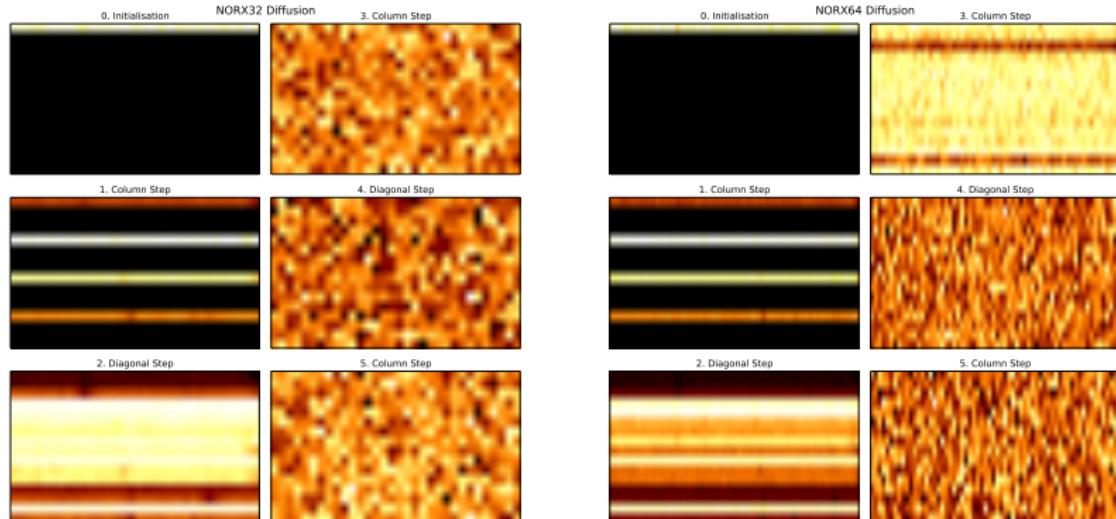


Figure: Diffusion Visualisation of  $F^R$ .

## NODE – The (NO)RX (D)ifferential Search (E)ngine\*

- ▶ Framework for automatic search of differentials in  $F^R$
- ▶ Examined differential propagation through  $H$
- ▶ Uses constraint / SAT solvers (STP, Boolector, CryptoMiniSat)
- ▶ Available on GitHub: <https://github.com/norx/NODE>
- ▶ Best differential trails in  $F^4$  (full state):  
 $2^{-584}$  (32-bit) and  $2^{-836}$  (64-bit)
- ▶ Differential trail bounds for  $F$  (init., diffs in nonce only):  
 $< 2^{-60}$  (32-bit) and  $< 2^{-53}$  (64-bit)
- ▶ Variant of NODE allowed us to break Wheesht and McMambo, two other CAESAR candidates

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## Algebraic Properties of G

| W      | #polynomials by degree |   |     |   |   |     | #monomials |     |     |        |
|--------|------------------------|---|-----|---|---|-----|------------|-----|-----|--------|
|        | 3                      | 4 | 5   | 6 | 7 | 8   | min        | max | avg | median |
| 32-bit | 2                      | 6 | 58  | 2 | 8 | 52  | 12         | 489 | 242 | 49.5   |
| 64-bit | 2                      | 6 | 122 | 2 | 8 | 116 | 12         | 489 | 253 | 49.5   |

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- ▶ High polynomial degree + big number of monomials + large state size: should increase difficulty to mount algebraic attacks

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We also examined weak states, fixed points, rotational properties, ...

## Algebraic Properties of G

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# NORX vs AES-GCM

|                       | NORX                                   | AES-GCM                           |
|-----------------------|--|-----------------------------------|
| High performance      | yes (on many platforms)                | depends (high with AES-NI)        |
| High key agility      | yes                                    | no                                |
| Timing resistance     | yes                                    | no (bit-slicing, AES-NI required) |
| Misuse resistance     | $A+N / LCP+X$ (exposes $P \oplus P'$ ) | no (exposes $K$ )                 |
| Parallelisation       | yes                                    | yes                               |
| Extensibility         | yes (sessions, secret msg. nr., etc.)  | no                                |
| Simple implementation | yes                                    | no                                |

# Conclusion

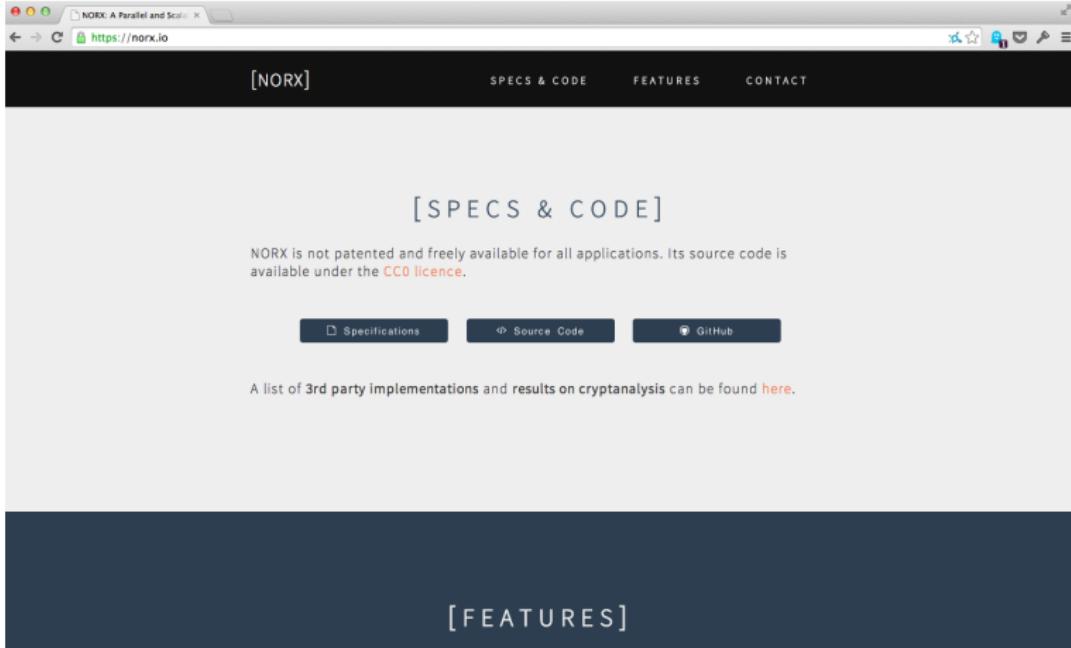
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- ▶ Fast on a broad range of architectures
- ▶ Resistance vs timing attacks in hw & sw (no Int. add. & no SBoxes)
- ▶ Our analysis found no security flaws
- ▶ Attacks only on reduced versions / single components
- ▶ NORX permutation probably a little bit weaker than ChaCha's
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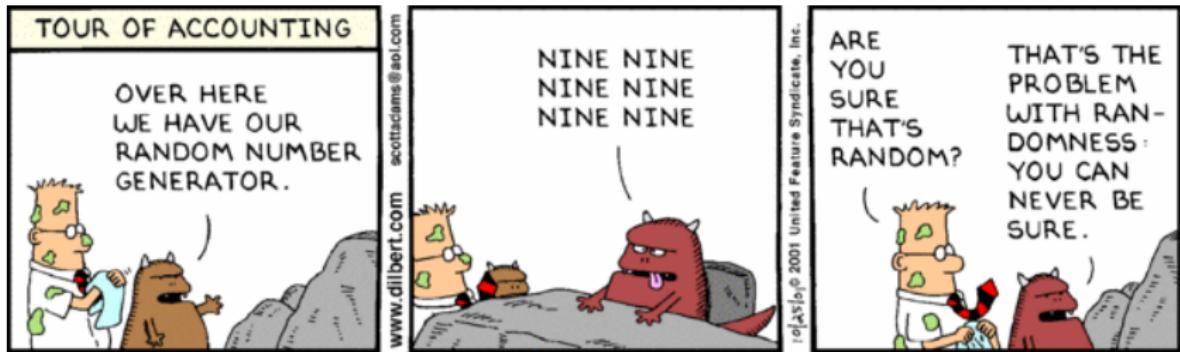
# Further Information



The screenshot shows a web browser window for the NORX website at <https://norx.io>. The page has a dark header with the title '[NORX]' and navigation links for 'SPECS & CODE', 'FEATURES', and 'CONTACT'. Below the header, a large section is titled '[ SPECS & CODE ]'. It contains text about the NORX algorithm being not patented and freely available under the CC0 licence, along with links for 'Specifications', 'Source Code', and 'GitHub'. A note mentions 3rd party implementations and cryptanalysis results. The bottom section is titled '[ FEATURES ]'.

<https://norx.io>

- ▶ J-P. Aumasson, P. Jovanovic, and S. Neves, NORX – A First Round Candidate in CAESAR
- ▶ J-P. Aumasson, P. Jovanovic, and S. Neves, NORX: Parallel and Scalable AEAD, European Symposium on Research in Computer Security (ESORICS 2014), Wroclaw, Poland
- ▶ J-P. Aumasson, P. Jovanovic, and S. Neves, Analysis of NORX, Cryptology ePrint Archive: Report 2014/317
- ▶ P. Jovanovic, A. Luykx, and B. Mennink, Beyond  $2^{c/2}$  Security in Sponge-Based Authenticated Encryption Modes, Cryptology ePrint Archive: Report 2014/373



Comic by <http://dilbert.com>

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